

SCJ2013 Data Structure & Algorithms

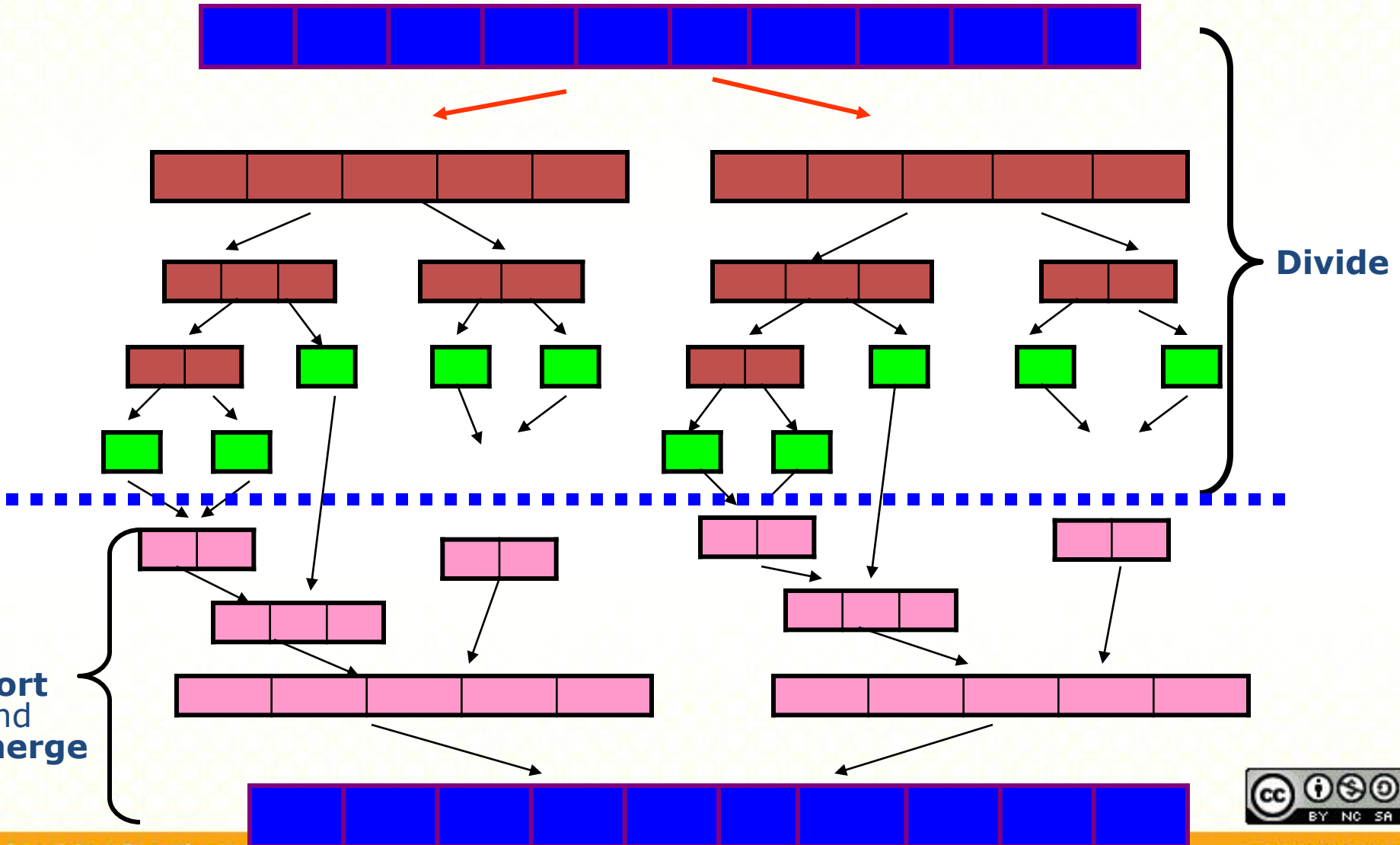
Merge Sort

Nor Bahiah Hj Ahmad & Dayang
Norhayati A. Jawawi

Merge Sort

- Merge Sort applies **divide** and **conquer** strategy.
- Three main steps in Merge Sort algorithm:
 - **Divide** an array into halves until only one piece left
 - **Sort** each half
 - **Merge** the sorted halves into one sorted array
- A recursive sorting algorithm
- **Performance** is independent of the **initial order** of the array items

Merge Sort Operation



Merge Sort Implementation


- **MergeSort ()** function
 - A **Recursive** function that **divide** the array into pieces until each piece contain only one item.
 - The small pieces are **merged** into larger sorted pieces until one sorted array is achieved.

mergeSort () function


```

void mergeSort(DataType theArray[],int first,int last)
{ if (first < last)
  { // sort each half
    int mid = (first + last)/2;
    // index of midpoint
    // sort left half theArray[first..mid]
    mergesort(theArray, first, mid);
    // sort right half theArray[mid+1..last]
    mergesort(theArray, mid+1, last);
    // merge the two halves
    merge(theArray, first, mid, last);
  } // end if
} // end mergesort
  
```

divide the
array into
pieces

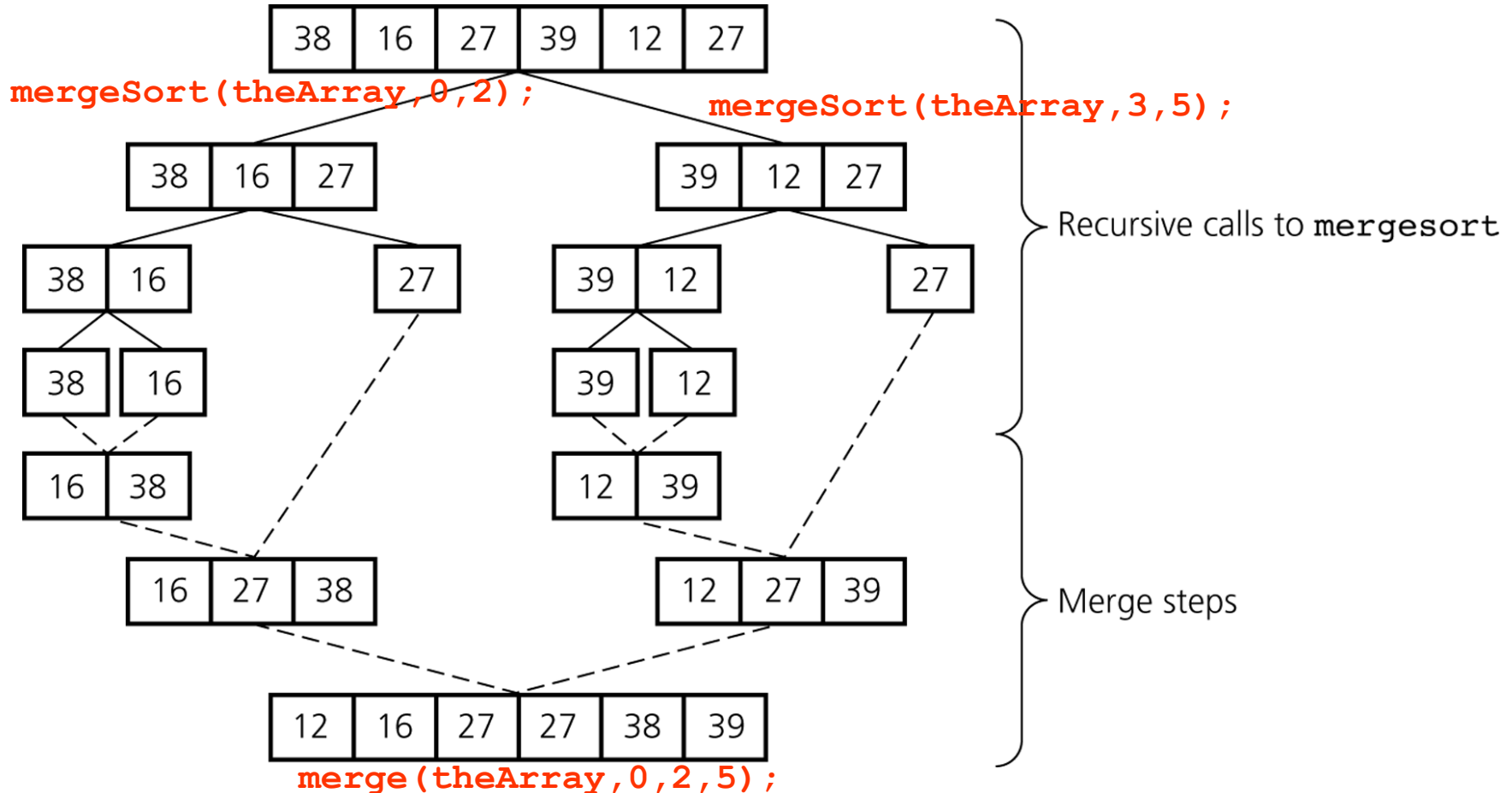


small pieces
are **merged**



mergeSort [38 16 27 39 12 27]

`mergeSort (theArray, 0, 5) ;`



A mergesort of an array of six integers

Merge Sort Implementation

- **Merge ()** function
 - **Compares** an item into one half of the array with item in the other half of the array and
 - **Moves** the smaller item into temporary array.
 - Then, the **remaining** items are simply **moved** to the temporary array.
 - The temporary array is copied back into the original array.

merge () function

```
const int MAX_SIZE = maxNmbrItemInArray;
void merge(DataType theArray[],
           int first, int mid, int last)
{   DataType tempArray[MAX_SIZE]; // temp array
    int first1 = first; // first subarray begin
    int last1  = mid;   // end of first subarray
    int first2 = mid + 1; // secnd subarry begin
    int last2  = last;   // end of secnd subarry
```

Duplicate
the positions

```
// while both subarrays are not empty,
// copy the smaller item into the
// temporary array
    int index = first1;
// next available location in tempArray
for (; (first1 <= last1) && (first2 <=
    last2); ++index)
{if (theArray[first1] < theArray[first2])
    { tempArray[index] = theArray[first1];
      ++first1;      }
  else
    { tempArray[index] = theArray[first2];
      ++first2;      }
} // end if
```

Moves the smaller
item into
temporary array

merge () function

```
for (; first1 <= last1;
    ++first1, ++index)
    tempArray[index] =
        theArray[first1];
// finish off the second
// subarray, if necessary

for (; first2 <= last2;
    ++first2, ++index)
    tempArray[index] =
        theArray[first2];
```

move the **remaining** items to the temporary array

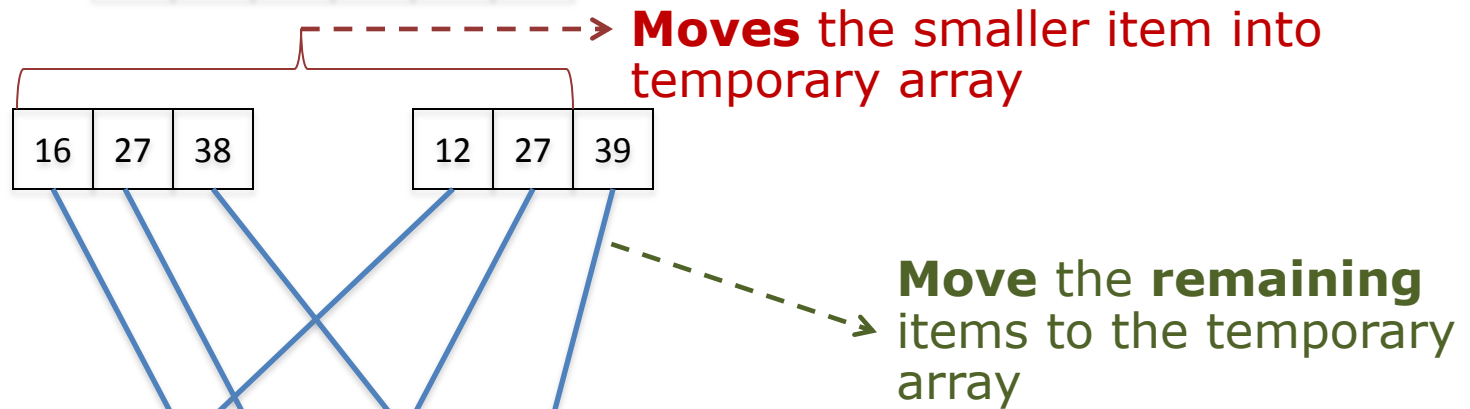
```
// copy the result back
// into the original
// array
for (index = first;
    index <= last; ++index)
    theArray[index] =
        tempArray[index];
} // end merge function
```

The temporary array is copied back into the original array

Merge Sort Operation

The array :

38	16	27	39	12	27
----	----	----	----	----	----



Temporary array

tempArray :

12	16	27	27	38	39
----	----	----	----	----	----

The array :

12	16	27	27	38	39
----	----	----	----	----	----

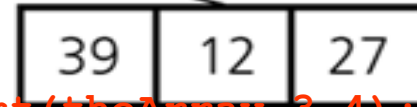
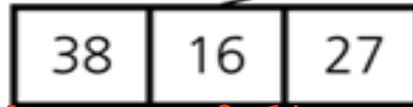
The temporary array is copied back into the original array

mergeSort [38 16 27 39 12 27]

1. mergeSort (theArray, 0, 5);

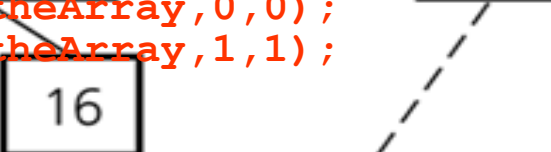


2. mergeSort (theArray, 0, 2); 9. mergeSort (theArray, 3, 5);



3. mergeSort (theArray, 0, 1);

10. mergeSort (theArray, 3, 4);

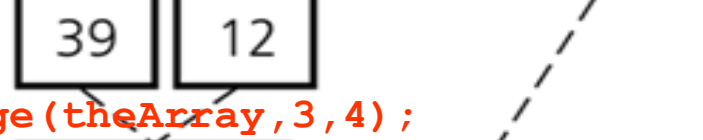


4. mergeSort (theArray, 0, 0);

11. mergeSort (theArray, 3, 3);

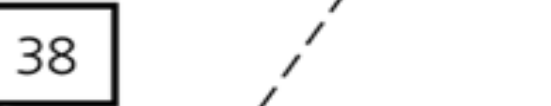
5. mergeSort (theArray, 1, 1);

12. mergeSort (theArray, 4, 4);



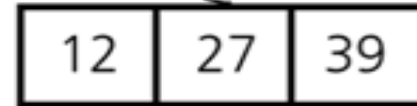
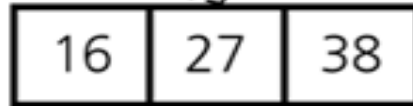
6. merge (theArray, 1, 1);

13. merge (theArray, 3, 4);

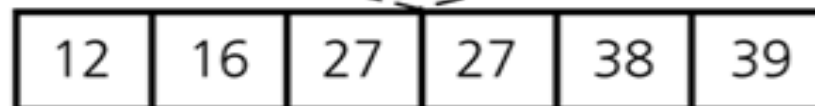


8. merge (theArray, 1, 1);

15. merge (theArray, 3, 4);



16. merge (theArray, 0, 5);



mergeSort [38 16 27 39 12 27]

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Unsorted data [38 16 27 39 12 27]

Content of divided sublist with first=0 & last=5 [38 16 27 39 12 27]

Content of divided sublist with first=0 & last=2 [38 16 27]

Content of divided sublist with first=0 & last=1 [38 16]

Content of divided sublist with first=0 & last=0 [38]

Content of divided sublist with first=1 & last=1 [16]

Content of merged list with first=0 & last=1 [16 38]

Content of divided sublist with first=2 & last=2 [27]

Content of merged list with first=0 & last=2 [16 27 38]

Content of divided sublist with first=3 & last=5 [39 12 27]

Content of divided sublist with first=3 & last=4 [39 12]

Content of divided sublist with first=3 & last=3 [39]

Content of divided sublist with first=4 & last=4 [12]

Content of merged list with first=3 & last=4 [12 39]

Content of divided sublist with first=5 & last=5 [27]

Content of merged list with first=3 & last=5 [12 27 39]

Content of merged list with first=0 & last=5 [12 16 27 27 38 39]

Sorted data [12 16 27 27 38 39]

Press any key to continue . . .



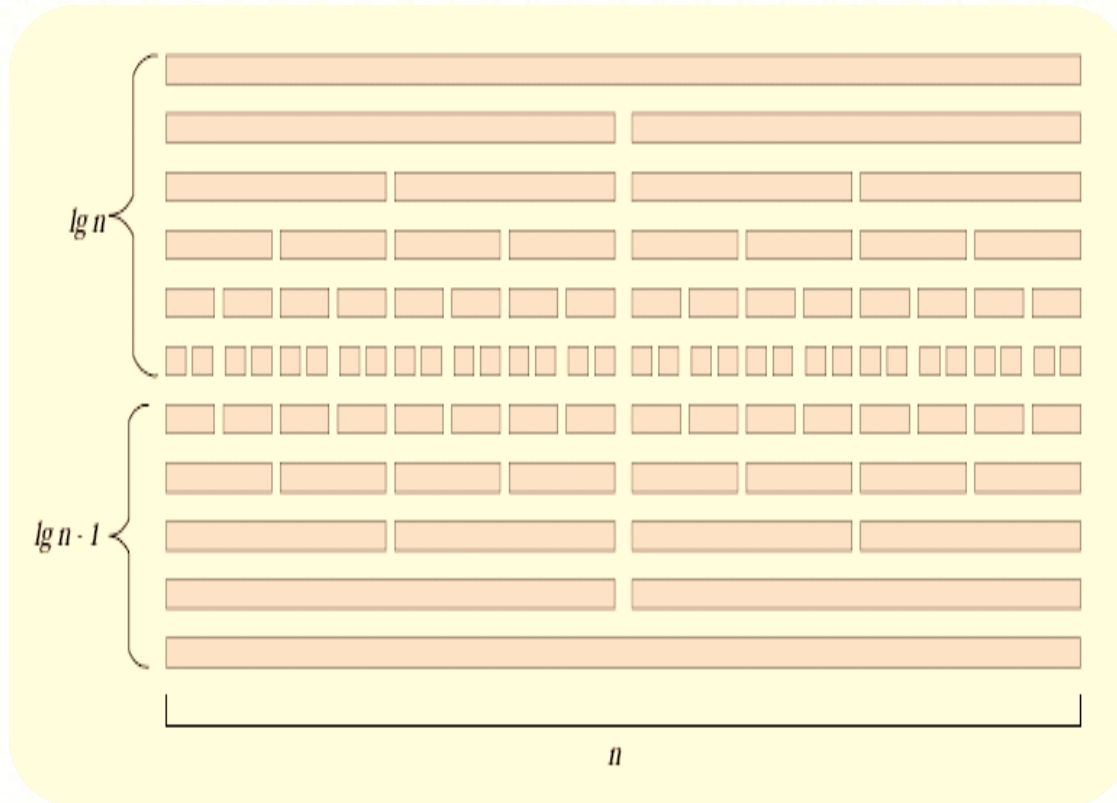
Merge Sort Analysis

- The list is **always divided into two balanced** list (or almost balanced for odd size of list)
- The number of calls to repeatedly divide the list until there is one item left in the list is:

$$n + 2 \frac{n}{2} + 4 \frac{n}{4} + 8 \frac{n}{8} + 16 \frac{n}{16} + \dots x \frac{n}{x}$$

- Assuming that the left segment and the right segment of the list have the equal size (or almost equal size), then $x \approx \lg n$.
The number of iteration is approximately $n \lg n$

Merge Sort Analysis



The same number of repetition is needed to sort and merge the list.

Thus, as a whole the number of steps needed to sort data using merge sort is $2n \lg n$, which is

$$O(n \lg n).$$

Mergesort

- Analysis
 - Worst case: $O(n * \log_2 n)$
 - Average case: $O(n * \log_2 n)$
 - Performance is independent of the initial order of the array items
- Advantage
 - Mergesort is an extremely fast algorithm
- Disadvantage
 - Mergesort requires a second array (temporary array) as large as the original array

References

1. Frank M. Carano, Janet J Prichard. “*Data Abstraction and problem solving with C++ Walls and Mirrors*. 5th edition (2007). Addison Wesley.
2. Nor Bahiah et al. *Struktur data & algoritma menggunakan C++*. Penerbit UTM, 2005.